

# QUASIRANDOM-FORCING TOURNAMENTS

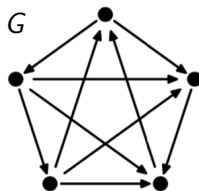
Fiona Skerman (Uppsala University)

joint work with Robert Hancock, Adam Kabela, Dan Král',  
Taísa Martins, Roberto Parente and Jan Volec

# We recall tournaments

## Definitions

A *tournament* is a directed graph having precisely one arc between each pair of its nodes.



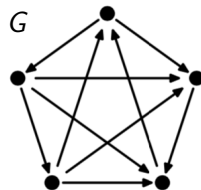
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$d(H, G)$  = probability that  $|H|$  randomly chosen vertices of  $G$  induce  $H$

$$d(H, G) = \frac{n(H, G)}{\binom{|G|}{|H|}} = \frac{8}{\binom{5}{3}}$$



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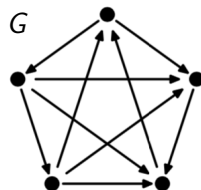
$d(H, G)$  = probability that  $|H|$  randomly chosen vertices of  $G$  induce  $H$

$d^*(H, G)$  = probability that an ordered set of  $|H|$  randomly chosen vertices of  $G$  induces a labelled copy of  $H$

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$$d^*(H, G) = \frac{n^*(H, G)}{|G|_{|H|}} = \frac{8}{5 \cdot 4 \cdot 3}$$

$$d^*(H, G) = \frac{|\text{Aut}(H)|}{|H|!} d(H, G)$$



# Quasirandomness of tournaments

**Definitions** Chung-Graham '91 (Chung-Graph-Wilson '89, Thomason 87')

Let  $(G_n)$  be sequence of tournaments (such that  $|G_n| \rightarrow \infty$  as  $n \rightarrow \infty$ ).

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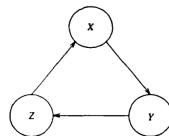
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**Thm (Chung-Graham '91)**

Let  $h \geq 4$  and define  $\mathcal{H}_h$  to be the set of tournaments on  $h$  nodes. Then  $\mathcal{H}_h$  is quasirandom-forcing.



$T^*$

FIGURE 1

Chung-Graham '91

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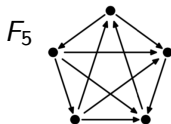
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A tournament  $H$  is *quasirandom-forcing* if

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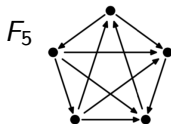
# Which tournaments on $h$ nodes are quasirandom-forcing?

- For  $h \geq 4$ , every transitive tournament is quasirandom-forcing (Coregliano-Razborov '17, Lovasz '93)
- For  $h = 5$ ,  $F_5$  is quasirandom-forcing, others not except perhaps  $H_5$ . (Coregliano, Parente and Sato '19).
- For  $h \geq 7$ , the only quasirandom-forcing is the transitive (Bucić, Long, Shapira and Sudakov, '20+). Also local-forcing.
- For  $h = 6$ , the only quasirandom-forcing is the transitive.  $H_5$  not. (Hancock, K., Král', Martins, Parente, Skerman and Volec, '20+).



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**Summary:**  $H$  is quasirandom-forcing iff  $H$  is  $F_5$  or  $H$  is the transitive tournament on  $h \geq 4$  nodes.

cf. Goodman '59, Beineke, Harary '65.

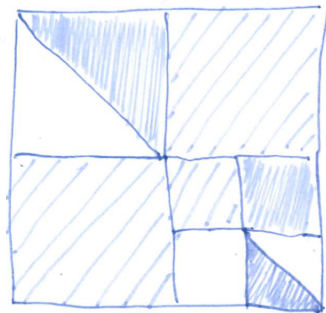
# General Arguments

## Proposition (Bucić, Long, Shapira and Sudakov '20+)

Let  $H$  be a non-transitive tournament on  $n \geq 7$  nodes. Then  $H$  is not quasirandom-forcing.

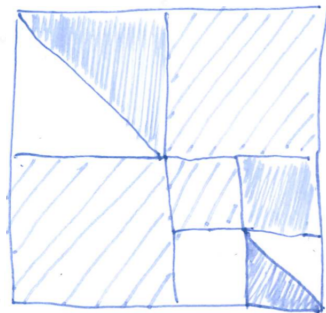
## Lemma (Hancock, Kabela, Král', Martins, Parente, S. and Volec '20+)

Let  $H$  be a non-transitive tournament on 6 nodes. If  $H$  **contains twins** or has a **non-trivial automorphism group** or is **not strongly connected** then  $H$  is not quasirandom-forcing.

$$W : [0, 1]^2 \rightarrow [0, 1] \text{ with } W(x, y) = 1 - W(y, x), \forall x, y.$$


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$$W\text{-random graph } S = \mathbb{G}(h, W).$$

sample  $x_1, \dots, x_h \in^u [0, 1]$ , for  $i < j$ :

 $\vec{ij} \in E(S)$  with probability  $W(x_i, x_j)$  $\vec{ji} \in E(S)$  otherwise.

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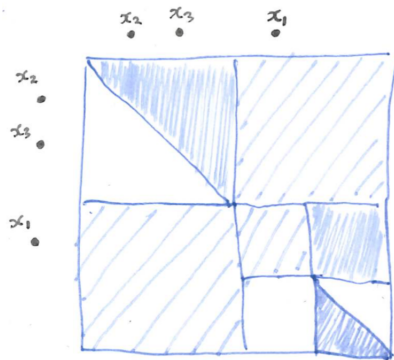
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$$W(x_1, x_2) = \frac{1}{2}$$

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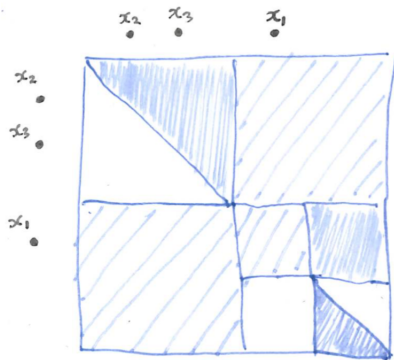
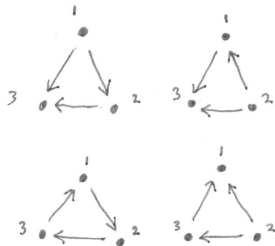
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$$\begin{aligned} W(x_1, x_2) &= \frac{1}{2} \\ W(x_1, x_3) &= \frac{1}{2} \\ W(x_2, x_3) &= 1 \end{aligned}$$



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labelled density

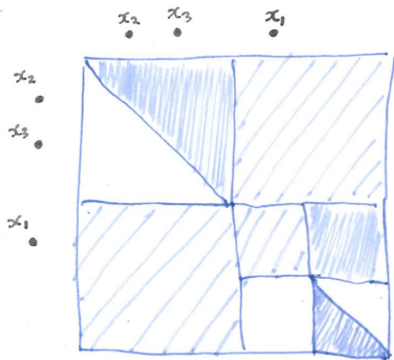
$$d^*(H, W) = \mathbb{P}(S =^{\text{labelled}} H)$$

$$= \int_{[0,1]^h} \prod_{\vec{ij} \in E(H)} W(x_i, x_j) dx_1 \cdots dx_h.$$

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6 / 10

# General Arguments

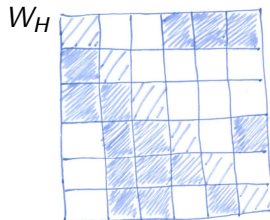
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$H$

0 0 0 1 1 1  
1 0 0 0 0 0  
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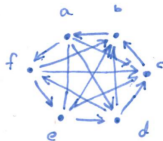
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Pf (BLSS):  $d^*(H, W_H) \geq h^{-h}$ . For  $h \geq 7$ ,  $h^{-h} \geq 2^{-\binom{h}{2}}$   $\square$ .

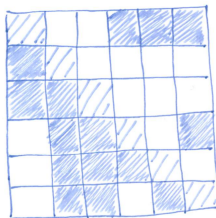
$$6^{-6} < 2^{-15} < 2 \times 6^{-6}$$

$H$



○	○	○			
	○	○	○	○	○
		○	○	○	○
○			○	○	
○				○	○
○			○		○

$W_H$



## General Arguments

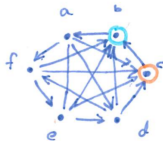
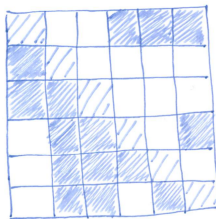
55 + 1

Lemma (Hancock, Kabela, Král', Martins, Parente, S. and Volec, 2020+)

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Let  $N^+(x)$  denote out-neighbours of  $x$ .

Nodes  $u$  and  $v$  are *twins* if  $N^+(x) \setminus \{y\} = N^+(y) \setminus \{x\}$ .

 $H$  $W_H$ 

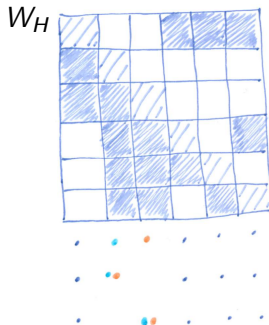
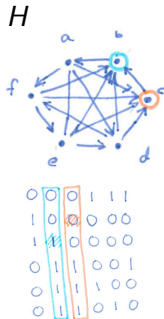
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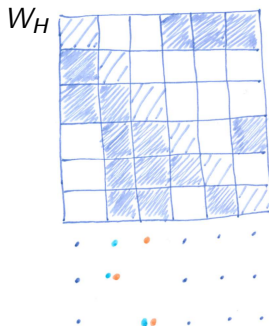
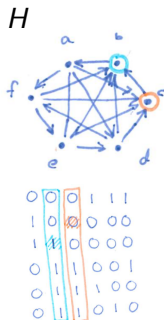
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Pf:  $d^*(H, W_H) \geq h^{-h}(1 + \frac{1}{2} + \frac{1}{2})$ . For  $h = 6$ ,  $2h^{-h} \geq 2^{-\binom{h}{2}}$   $\square$ .



# General Arguments

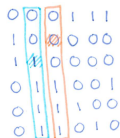
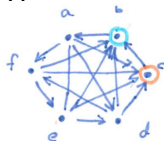
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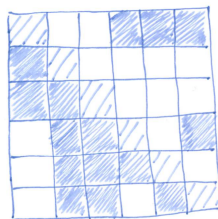
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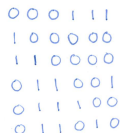
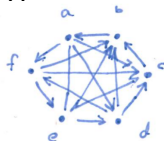
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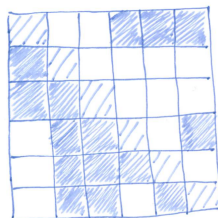
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# General Arguments

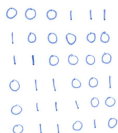
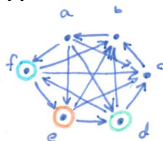
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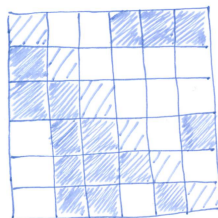
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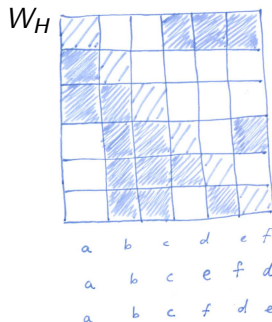
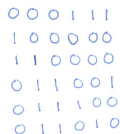
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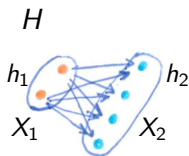
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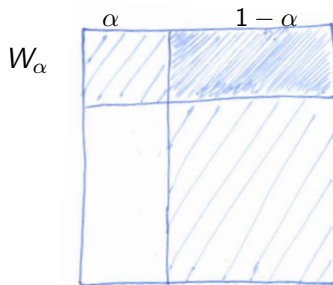
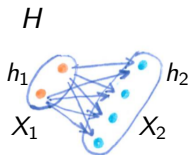
$$18 + 1$$

If  $H$  is **not strongly connected** then  $H$  is not quasirandom-forcing.


$$W_\alpha$$

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## General Arguments

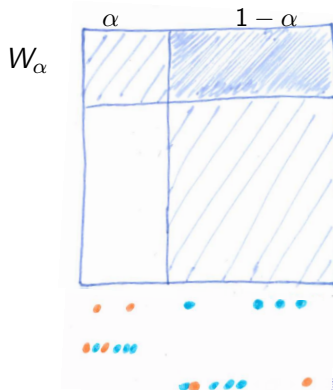
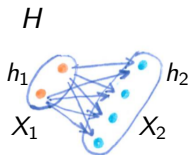
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Lemma (Hancock, Kabela, Král', Martins, Parente, S. and Volec, 2020+)

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$$\text{Pf: } d^*(H, W_\alpha) \geq \alpha^{h_1} (1-\alpha)^{h_2} 2^{-\binom{h_1}{2} - \binom{h_2}{2}} + \alpha^{h_2} 2^{-\binom{h}{2}} + (1-\alpha)^{h_2} 2^{-\binom{h}{2}}.$$



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14 + 1

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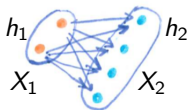
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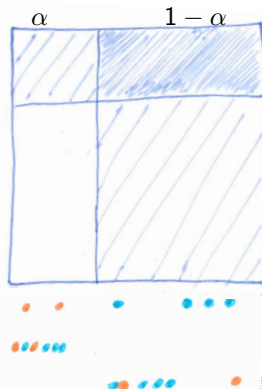
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$$\Rightarrow \exists \alpha, d^*(H, W_\alpha) > 2^{-\binom{h}{2}} \square.$$

$H$



$W_\alpha$

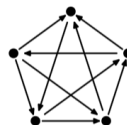
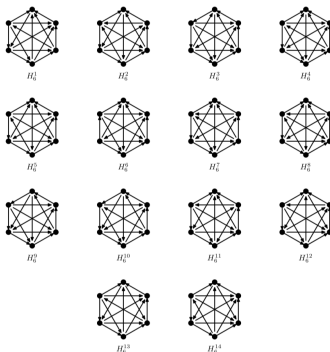


# General Arguments

14 + 1

Lemma (Hancock, Kabela, Král', Martins, Parente, S. and Volec '20+)

Let  $H$  be a non-transitive tournament on 6 nodes. If  $H$  contains **twins** or has a **non-trivial automorphism group** or is **not strongly connected** then  $H$  is not quasirandom-forcing.



## General Arguments

14 + 1

(A) not strongly  
connected(B) non-trivial auto-  
morphism group

(C) contains twins

A	B	C	D	E	Tournament
•	•	•			00000,0000,000,01,0
•		•			00010,0000,000,00,0
		•			00011,0000,000,00,0
		•			00010,0001,000,00,0
		•	•		00010,0000,001,00,0
		•			00010,0000,000,01,0
		•			00010,0000,000,00,1
•		•			00000,0010,000,00,0
		•			00001,0010,000,00,0
•		•			00000,0011,000,00,0
•					00000,0010,001,00,0
•					00000,0010,000,01,0
•		•			00000,0010,000,00,1
•	•				00000,0011,001,00,0
•	•	•			00000,0000,010,00,0
•		•			00001,0000,010,00,0
•	•				00000,0001,010,00,0
•		•			00000,0000,011,00,0
•		•			00100,0000,000,00,0
•		•			00110,0000,000,00,0
		•			00111,0000,000,00,0
		•			00110,0001,000,00,0
		•			00110,0000,001,00,0
		•			00110,0000,000,01,0
		•			00110,0000,000,00,1
		•			00111,0000,001,00,0
•		•			00110,0001,001,00,0
•	•				00111,0000,000,01,0

 $H_6^1$ 

A	B	C	D	E	Tournament
		•			00110,0001,000,01,0
•					00100,0010,000,00,0
		•			00101,0010,000,00,0
		•			00100,0011,000,00,0
		•			00100,0010,001,00,0
		•			00100,0010,000,01,0
		•			00100,0010,000,00,1
	•				00101,0010,001,00,0
		•			00100,0011,001,00,0
		•			00100,0011,000,01,0
•	•				00110,0010,000,00,0
		•			00111,0010,000,00,0
		•			00111,0011,000,00,0
		•			00111,0010,001,00,0
•	•	•			00000,0100,000,00,0
•	•				00010,0100,000,00,0
	•	•			00011,0100,000,00,0
		•			00010,0101,000,00,0
	•				00010,0100,000,00,1
•	•	•			01000,0000,000,00,0
•	•				01000,0000,000,01,0
•					01010,0000,000,00,0
		•			01011,0000,000,00,0
		•			01010,0001,000,00,0
		•			01010,0000,001,00,0
		•			01010,0000,000,01,0
	•				01010,0000,000,00,1

 $H_6^2$  $H_6^3$  $H_6^4$  $H_6^5$  $H_6^6$  $H_6^7$  $H_6^8$  $H_6^9$  $H_6^{10}$  $H_6^{11}$  $H_6^{12}$  $H_6^{13}$  $H_6^{14}$

# Excluding the rest

14 + 1

We readily check the properties from the previous slide and consider the 14 remaining tournaments on 6 nodes plus 1 tournament on 5 nodes.

To show tournament  $H$  on  $h$  nodes is not quasirandom-forcing

- find tournament  $T$  with many copies of  $H$ .

$$n(H, T) > |T|^h 2^{-\binom{h}{2}} \Rightarrow d^*(H, W_T) > 2^{-\binom{h}{2}}$$

- find a step tournament by perturbing about  $1/2$ .

$$d^*(H, A_x) = f(x) \text{ and find } x \text{ such that } f(x) > 2^{-\binom{h}{2}}.$$

$$A_x = \begin{pmatrix} 1/2 & 1/2 - x \\ 1/2 + x & 1/2 \end{pmatrix},$$

$$B_x = \begin{pmatrix} 1/2 & 1/2 - x & 1/2 + x \\ 1/2 + x & 1/2 & 1/2 - x \\ 1/2 - x & 1/2 + x & 1/2 \end{pmatrix},$$

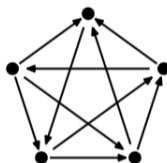
$$C_x = \begin{pmatrix} 1/2 & 1/2 - x & 1/2 + x & 1/2 - x \\ 1/2 + x & 1/2 & 1/2 - x & 1/2 - x \\ 1/2 - x & 1/2 + x & 1/2 & 1/2 - x \\ 1/2 + x & 1/2 + x & 1/2 + x & 1/2 \end{pmatrix}.$$

**9 + 1**

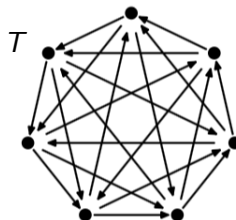
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$$n(H_5, T) = 21$$

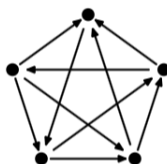


## 9

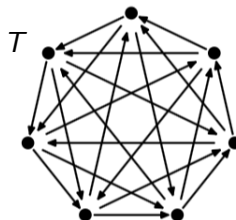
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$$n(H_5, T) = 21$$



Thank you for your attention.

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