

There are no first-order sentences with quantifier depth 4 and
an infinite spectrum

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Outline

- 1 First Order Spectra
- 2 Required theorems and constructions
- 3 How to play the Ehrenfeucht game
- 4 Proof sketch for $\frac{1}{2}$
- 5 How to win for Spoiler
- 6 Future research

1 First Order Spectra

Definition

Erdős–Rényi random graph model is a probabilistic space

$$G(N, p) = (\Omega_N, \mathcal{F}_N, \mathbb{P}_{N,p}),$$

where $N \in \mathbb{N}$, $0 \leq p \leq 1$,

$$\Omega_N = \{\mathcal{G} = (\mathcal{V}_N, \mathcal{E})\} -$$

set of all graphs with $\mathcal{V}_N = \{1, 2, \dots, N\}$,

$$\mathcal{F}_N = 2^{\Omega_N}, \quad \mathbb{P}_{N,p}(\mathcal{G}) = p^{e(\mathcal{G})} (1 - p)^{C_N^2 - e(\mathcal{G})}.$$

Definition

First-order graph property is a property defined by the first order formula with the following symbols:

- predicate symbols: $=, \sim$;
- logical symbols: $\rightarrow, \wedge, \vee, \neg, \dots$;
- variables: x, y, \dots ;
- quantifiers: \exists, \forall .

Definition

Quantifier depth of a first-order property Q is a minimal quantifier depth of a first-order formula that expresses Q .

Definition

For $p = p(N)$ the Zero-One Law holds, if for each first-order property L $P_{N,p}(L)$ tends to either 0 or 1.

Consider probabilities $p(N) = N^{-\alpha}$.

Theorem (J. Spencer, S. Shelah, 1988)

Let $p(N) = N^{-\alpha}$.

- Let α be a positive irrational. Then for $p(N)$ a Zero-one law holds.
- Let α be a positive rational. If $\alpha > 2$ or $\alpha \in \left(1 + \frac{1}{l+1}, 1 + \frac{1}{l}\right)$ for some $l \in \mathbb{N}$, then for $p(N)$ the Zero-One Law holds. In all other cases the Zero-One Law for $p(N)$ does not hold.

Definition

For $p = p(N)$ the Zero-One k -Law holds, if for each first-order property L with quantifier depth $\leq k$ $P_{N,p}(L)$ tends to either 0 or 1.

Theorem (M. Zhukovskii, 2012)

Let $p(N) = N^{-\alpha}$. If $\alpha \in \left(0, \frac{1}{k-2}\right)$, then for $p(N)$ a Zero-One k -Law holds.
If $\alpha = \frac{1}{k-2}$, then for $p(N)$ the Zero-One k -Law does not hold.

Definition

We call a k -spectrum a set of all $\alpha \in (0, 1)$ s.t. for $p(N) = N^{-\alpha}$ the Zero-One k -Law does not hold.

Consider a FO property Q .

Definition

Spectrum of Q is a set of all $\alpha \in (0, 1)$ s.t. $\mathbf{P}_{N,p}(Q)$ does not tend to either 0 or 1 for $p(N) = N^{-\alpha}$.

A k -spectrum is clearly a union of all spectra of FO formulas Q with quantifier depth $\leq k$.

Example: simple formula spectrum

- Consider property $Q = \{G \text{ has an induced } C_4\} \Delta \{G \text{ has an induced } K_4\}$.
- By the Bollobás theorem, $p_1 = \frac{1}{N}$ and $p_2 = N^{-\frac{2}{3}}$ are threshold probabilities of respective parts
- Hence, the spectrum equals $\left\{\frac{2}{3}, 1\right\}$ (well, not really)

Infinity of spectrum

Theorem (J. Spencer, 1990)

There exists a FO property of depth 14 with an infinite spectrum.

Theorem (M. Zhukovskii, 2016)

There exists a FO property of depth 5 with an infinite spectrum. The 5-spectrum is finite.

Theorem (M. Zhukovskii, A. Matushkin, 2017)

The only possible limiting points of 4-spectrum are $1/2$ and $3/5$.

Theorem (A. Matushkin, M. Zhukovskii, Y.Y.)

Points $\frac{1}{2}$ and $\frac{3}{5}$ cannot be limiting in 4-spectrum. Therefore, 4-spectrum is finite.

2 Required theorems and constructions

We now give a criterion of validity of the zero-one k -law.

Theorem (A. Ehrenfeucht, 1960)

Random graph $G(N, p)$ obeys the zero-one k -law if and only if

$$\lim_{N, M \rightarrow \infty} \mathbb{P}_{N, M, p} \left(\left\{ (A, B) : \text{Duplicator has a winning strategy} \right. \right. \\ \left. \left. \text{in gameEHR}(A, B, k) \right\} \right) = 1,$$

где $\mathbb{P}_{N, M, p}$ is a product of measures $\mathbb{P}_{N, p}$ и $\mathbb{P}_{M, p}$.

Threshold probabilities of properties “contain a copy of the subgraph”

Let

$$\rho(G) = \frac{e(G)}{v(G)};$$
$$\rho^{\max}(G) = \max_{H \subseteq G} \rho(H).$$

Theorem (Ruciński A., Vince A, 1985)

Let $p_0(N) = N^{-1/\rho^{\max}(G)}$. If $p = o(p_0)$ then $G(N, p)$ a.a.s. does not contain a copy of G . If $p = o(p)$ then, on the contrary, $G(N, p)$ a.a.s. does contain a copy of G .

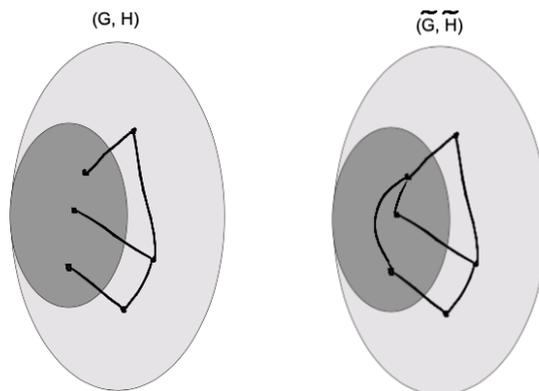
Definition

Let (G, H) and (\tilde{G}, \tilde{H}) , $G \subset H$, $\tilde{G} \subset \tilde{H}$ be two pairs of graphs. Let

$$V(G) = \{x_1, \dots, x_m\}, V(H) = \{x_1, \dots, x_l\},$$

$$V(\tilde{G}) = \{\tilde{x}_1, \dots, \tilde{x}_m\}, V(\tilde{H}) = \{\tilde{x}_1, \dots, \tilde{x}_l\}.$$

If $(x_i, x_j) \in E(G) \setminus E(H) \Rightarrow (\tilde{x}_i, \tilde{x}_j) \in E(\tilde{G}) \setminus E(\tilde{H})$, then \tilde{G} is called a (G, H) -extension of \tilde{H} .

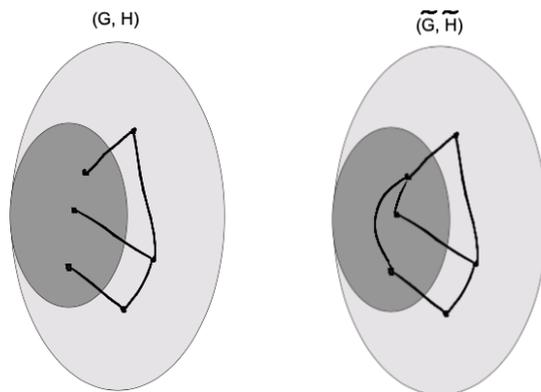


Fix $\alpha > 0$. Let

$$\begin{aligned}v(G, H) &= v(G) - v(H), e(G, H) = e(G) - e(H), \\f_\alpha(G, H) &= v(G, H) - \alpha e(G, H).\end{aligned}$$

$$f_\alpha(G, H) = v(G, H) - \alpha e(G, H).$$

Pair (G, H) is called α -safe, if $\forall S \quad (H \subset S \subseteq G \rightarrow f_\alpha(S, H) > 0)$.

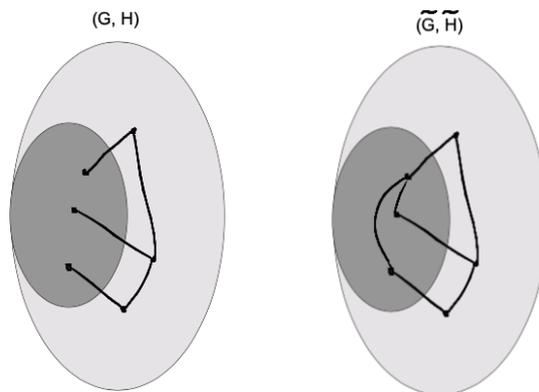


$\frac{1}{2}$ -safe

If the pair (G, H) is α -safe then a.a.s. there is (there are many) a (G, H) -extension of each subgraph \tilde{H} in the random graph $G(N, p)$.

$$f_\alpha(G, H) = v(G, H) - \alpha e(G, H).$$

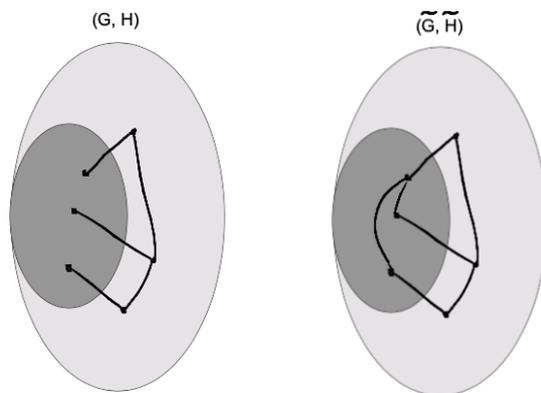
Pair (G, H) is called α -rigid, if $\forall S \quad (H \subseteq S \subset G \rightarrow f_\alpha(G, S) < 0)$.



$\frac{1}{3}$ -rigid

$$f_\alpha(G, H) = v(G, H) - \alpha e(G, H).$$

Pair (G, H) is called α -neutral, if $\forall S \quad (H \subset S \subset G \rightarrow f_\alpha(S, H) > 0)$ и $f_\alpha(G, H) = 0$



$\frac{3}{5}$ -neutral

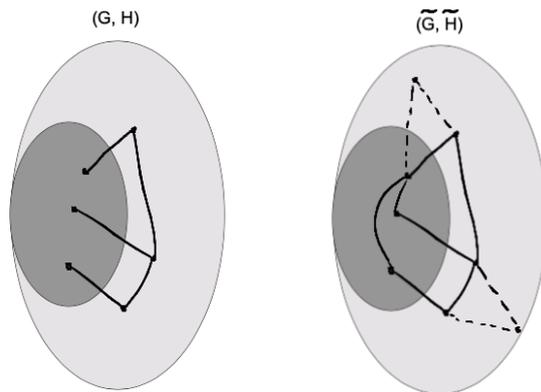
Maximal extensions

Question:

Can there be any α -neutral or α -rigid extensions in a random graph?

Answer:

Yes, there can be. But we can choose new vertices so that there are none.



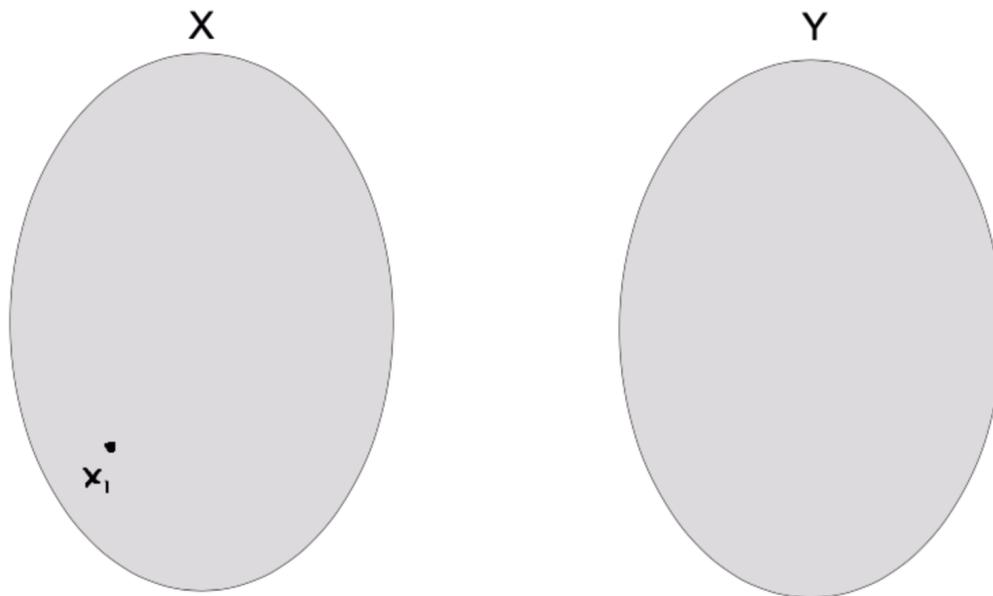
$$\alpha \in \left(\frac{1}{2}, \frac{1}{2} + \varepsilon\right)$$

We can choose a (G, H) -extension such that there are no outer vertices of degree ≥ 2

3 How to play the Ehrenfeucht game

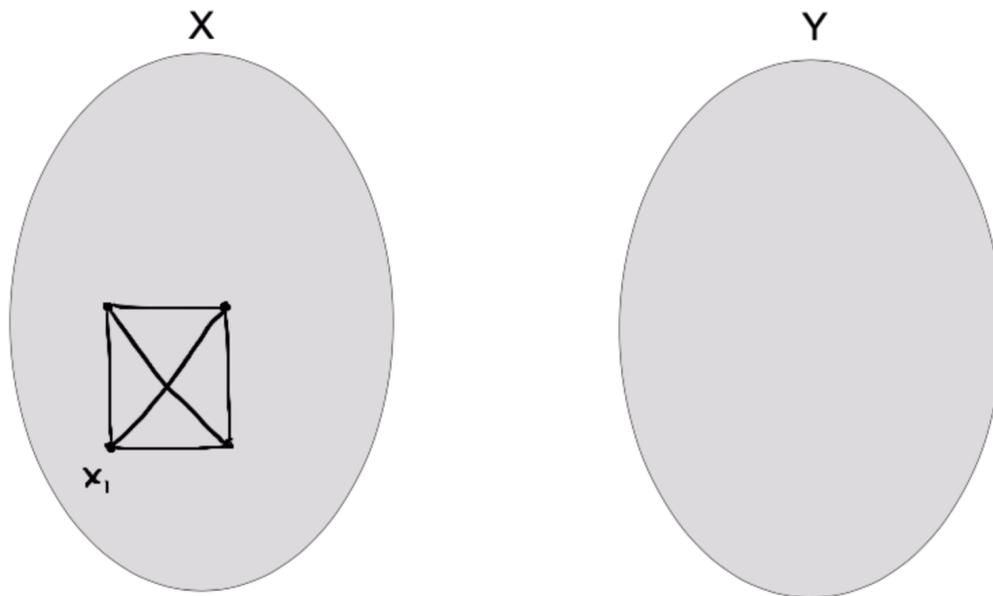
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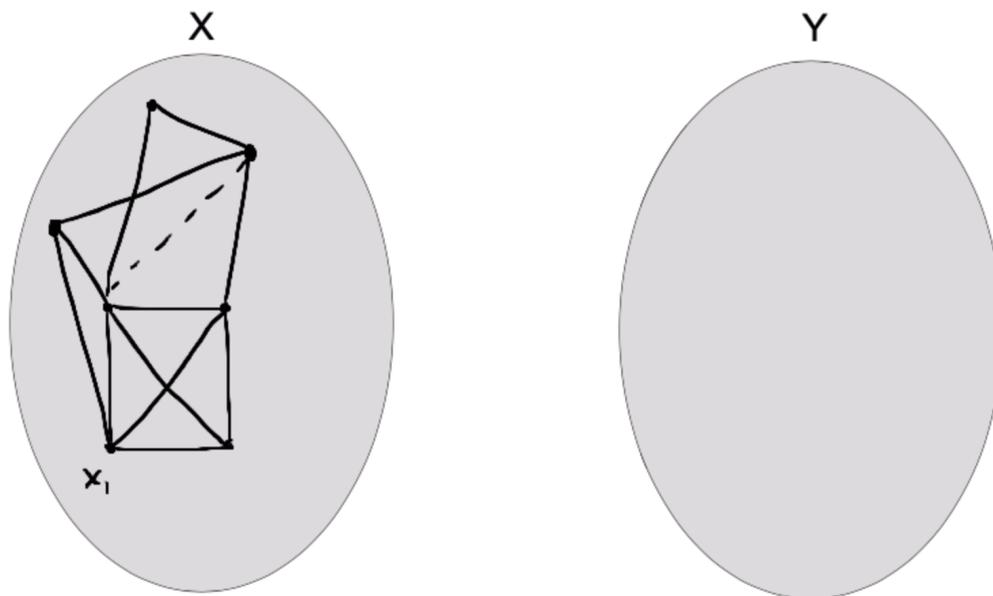
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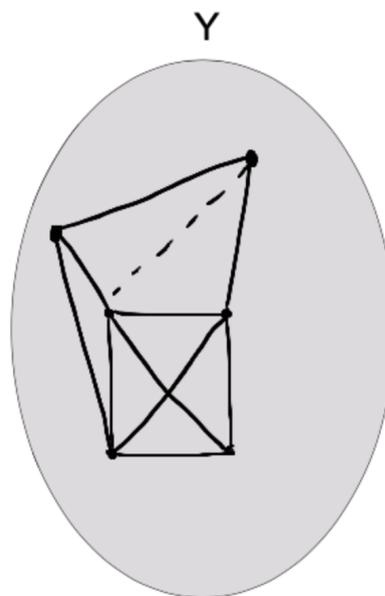
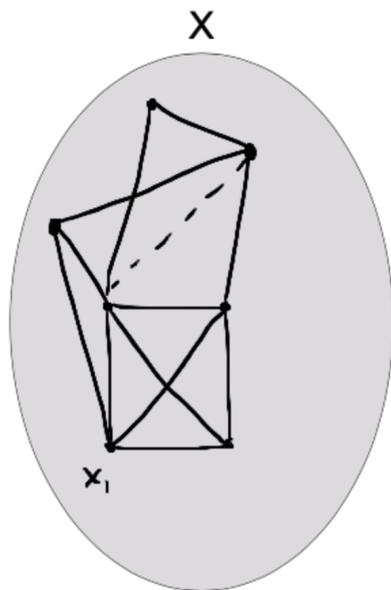
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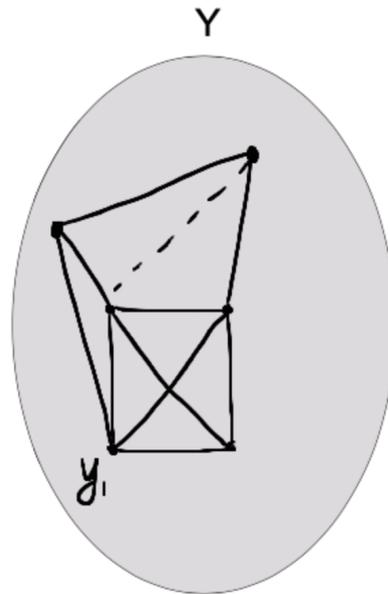
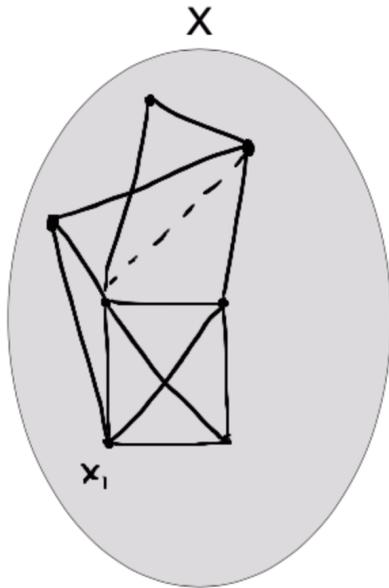
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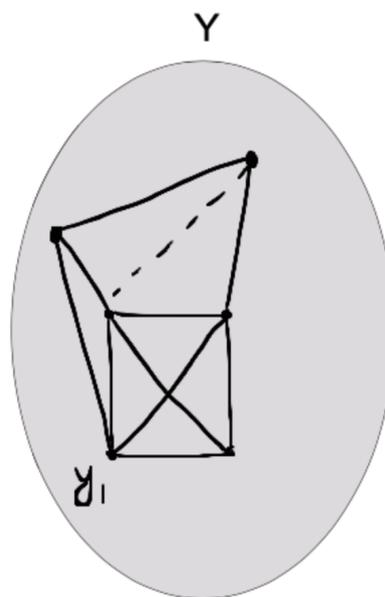
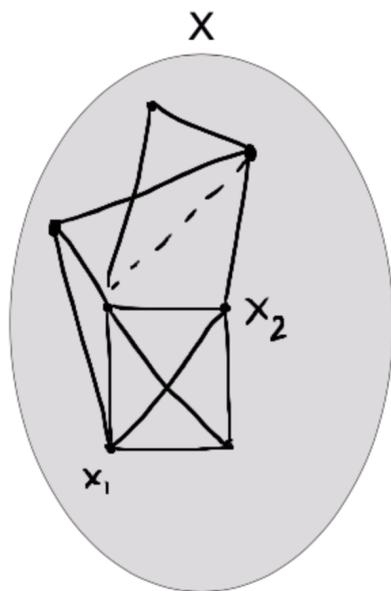
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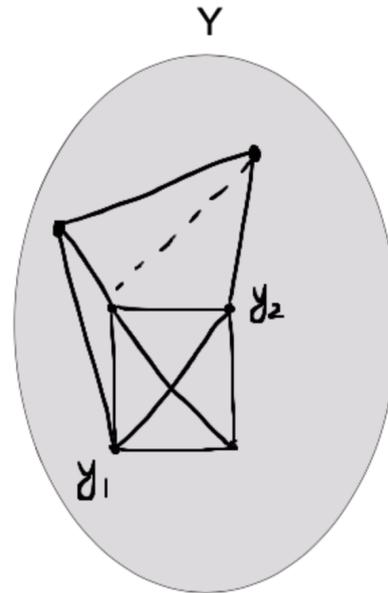
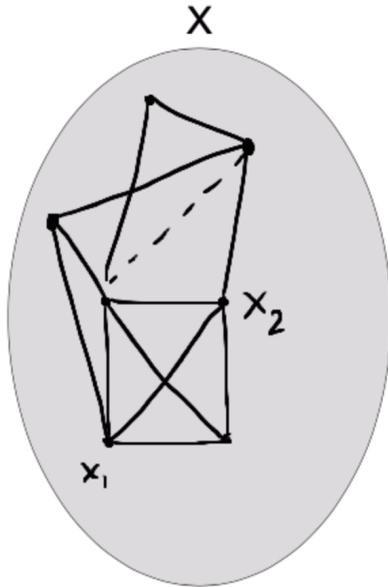
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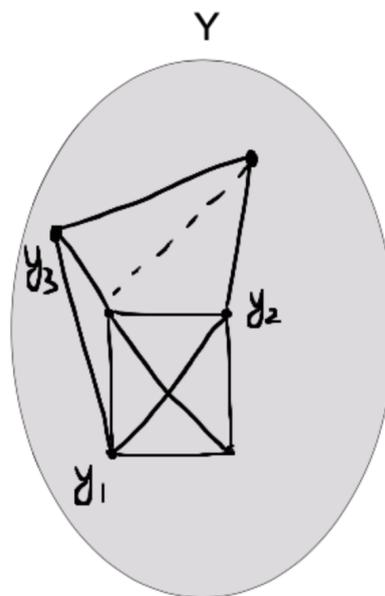
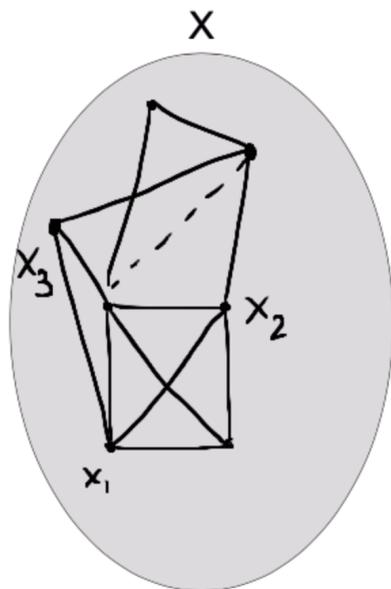
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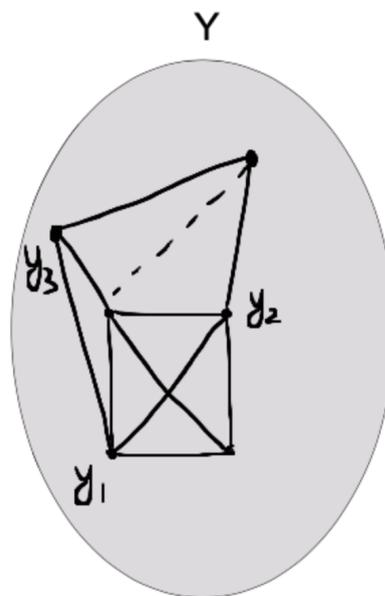
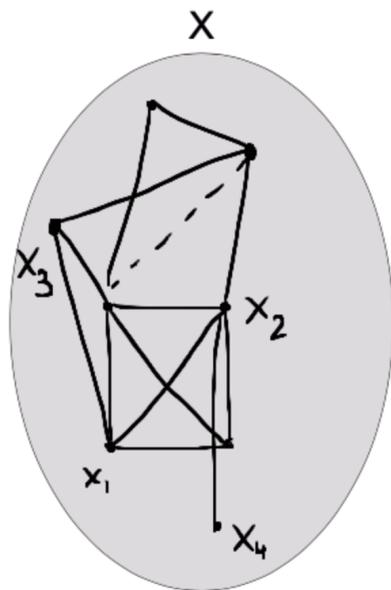
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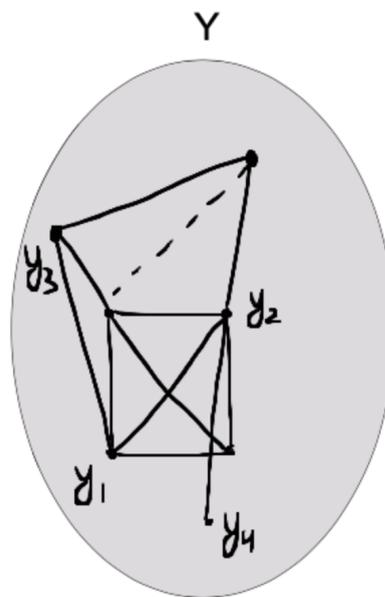
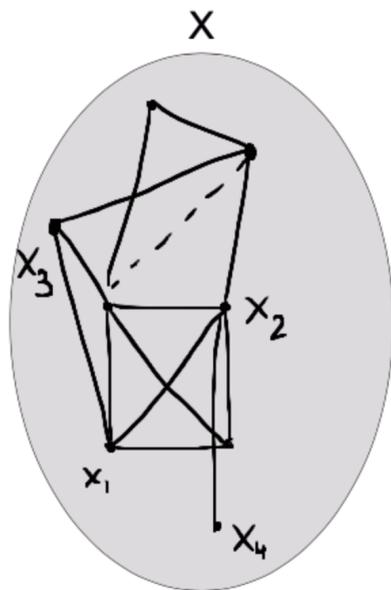
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How to play the Ehrenfeucht game

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4 Proof sketch for $\frac{1}{2}$

- We prove that there exists $\varepsilon > 0$ s.t. for each $\alpha \in (\frac{1}{2}, \frac{1}{2} + \varepsilon)$ a.a.s. in measure P_α Dupliterator wins.
- We give a strategy for Dupliterator to win a.a.s. in $\text{EHR}(X, Y, 4)$.
- Let $x_1 \in V(X)$ be the Spoiler's first step, let U be the maximal dense subgraph of X containing x_1 that "spoils" Dupliterator's play.
- Since $U \subseteq X$, we have a.a.s. $\rho^{\max}(U) < \frac{1}{\alpha}$.
- There is an "almost copy" W of graph U in Y , which is maximal.
- Dupliterator chooses $y_1 \in V(Y)$, which is the isomorphic image of x_1 .
- Similarities between U и W provide for the win of Dupliterator.

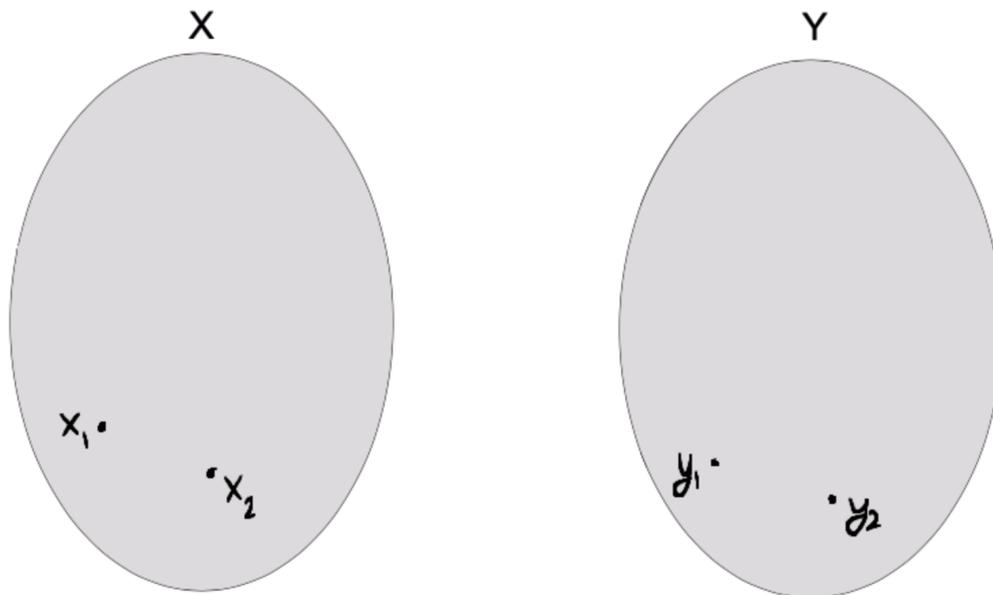
5 How to win for Spoiler

How to win for Spoiler

- Consider $\alpha = \frac{1}{2}$
- We prove that α is a limiting point in 5-spectrum
- We give a strategy for Spoiler to win in $\text{EHR}(X, Y, 5)$
- This proves the existence of the formula of depth 5 with an infinite spectrum.

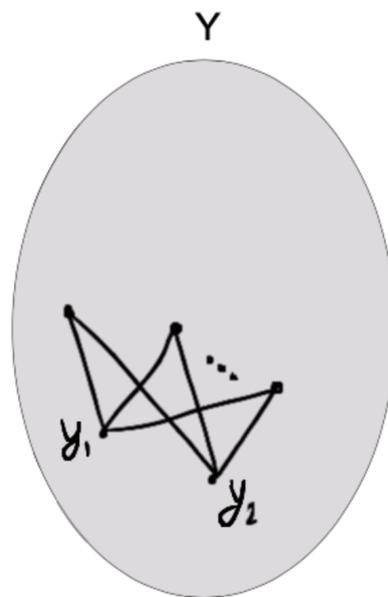
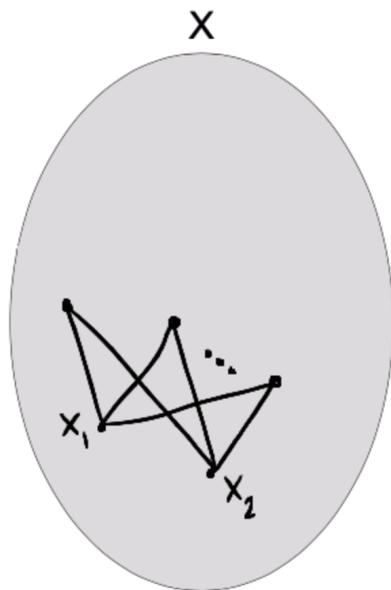
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$$\alpha \in \left(\frac{1}{2}, \frac{1}{2} + \varepsilon\right)$$



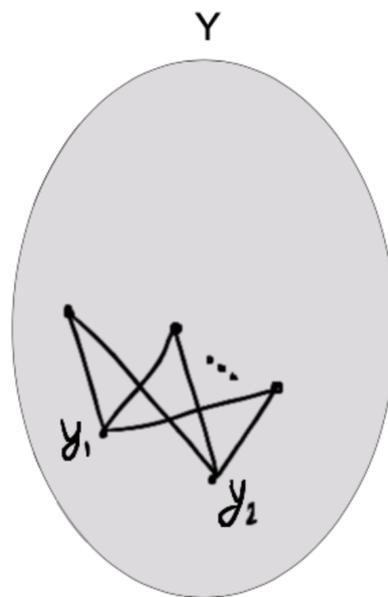
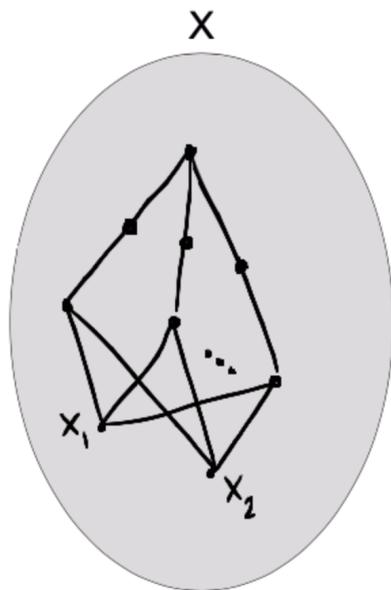
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How to win for Spoiler

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6 Future research

- Explain the structure of k -spectrum.
- Study the limiting points of k -spectrum for arbitrary k
- At least understand whether $\frac{1}{k-2}$ is limiting

Thank you for your attention!