

# A useful tool in combinatorics: Intersecting set-pair systems

Z. Füredi

(newest results with A. Gyárfás and Z. Király)

Rényi Institute of Mathematics, Budapest, Hungary

`z-furedi@illinois.edu`

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# Extended Abstract

A *cross intersecting set pair system* (SPS) of size  $m$ :

$(\{A_i\}_{i=1}^m, \{B_i\}_{i=1}^m)$  with  $A_i \cap B_i = \emptyset$  and  $A_i \cap B_j \neq \emptyset$ .

It is an important tool of extremal combinatorics. Bollobás' classical result states that  $m \leq \binom{a+b}{a}$  if  $|A_i| \leq a$  and  $|B_i| \leq b$  for each  $i$ .

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**1-cross intersecting** set pair system:  $|A_i \cap B_j| = 1$  for all  $i \neq j$ .

We show connections to **perfect graphs**, clique partitions of graphs, and **finite geometries**. The max size of a **1-cross intersecting SPS** is

- at least  $5^{n/2}$  for  $n$  even,  $a = b = n$ ,
- equal to  $(\lfloor \frac{n}{2} \rfloor + 1)(\lceil \frac{n}{2} \rceil + 1)$  if  $a = 2$  and  $b = n \geq 4$ ,
- at most  $|\cup_{i=1}^m A_i|$ ,
- asymptotically  $n^2$  if  $\{A_i\}$  is a **linear hypergraph**  
( $|A_i \cap A_j| \leq 1$  for  $i \neq j$ ).

# Some standard notation

$$[n] := \{1, 2, \dots, n\}$$

$$\binom{S}{k} := \text{set of } k\text{-sets, } 2^S \text{ power set}$$

$\deg_G(x)$  degree of vertex  $x$  of graph  $G = (V, E)$

$N_G(x) \subset V$ , neighborhood

$T \subseteq V(\mathcal{H})$  is a **cover** (transversal) of the hypergraph  $\mathcal{H} = (V, \mathcal{E})$   
if  $T \cap e \neq \emptyset \quad \forall e \in \mathcal{E}$ .

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$\mathcal{M} \subseteq \mathcal{E}$  is a **matching** (parallel edges) if  $e \cap e' = \emptyset \quad \forall e \neq e' \in \mathcal{M}$ .

$\nu(\mathcal{H}) := \max |\mathcal{M}|$ . (matching number of  $\mathcal{H}$ ).

$\nu = 1 \iff \mathcal{H}$  is **intersecting hpgr**.

# Cross intersecting set pair systems

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Theorem (Bollobás 1965)

$(A_1, B_1), (A_2, B_2), \dots, (A_m, B_m)$  set pairs with

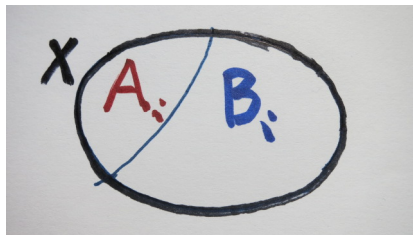
$|A_i| \leq a$ ,  $|B_i| \leq b$  and

$A_i \cap B_i = \emptyset$ , and

$A_i \cap B_j \neq \emptyset$  for all  $i \neq j$ , then  $m \leq \binom{a+b}{a}$ .

Best possible:

$A := \binom{X}{a}$ ,  $B := \binom{X}{b}$ ,



# Proof of Bollobás' theorem

Conj'd by Berge/ Ehrenfeucht, Mycielsky 1973. Proof by Jaeger, Payan 1971, Tarjan 1975, Katona 1974, (Alon / Frankl).

Given  $(\mathcal{A}, \mathcal{B})$ , a cross intersecting SPS.  $X := \cup \mathcal{A} \cup \mathcal{B}$ .

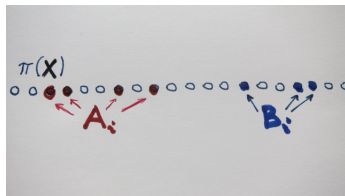


Figure: A type  $i$  permutation of  $X$ .

Call a permutation **type  $i$**  if  $A_i <_{\pi} B_i$ .

Note that type  $i \neq$  type  $j$ .



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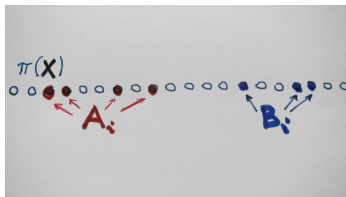


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Note that type  $i \neq$  type  $j$ .

$$\mathbb{P}(\text{prob}(\pi \text{ is of type } i)) = 1 / \binom{|A_i| + B_i|}{|A_i|}.$$

Therefore

$$\sum_i \frac{1}{\binom{|A_i| + B_i|}{|A_i|}} \leq 1. \quad \square$$

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2.

To investigate **1-cross intersecting** SPS. ( $|A_i \cap B_j| = 1$ )  
New results are joint with Gyárfás and Király, and related to one of my favorite structures, **finite affine and projective planes**

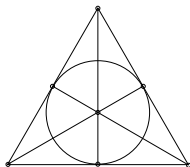


Figure: The Fano plane.

# The kernel of intersecting families

DEF: Suppose  $\mathcal{H}$  is intersecting.  $S$  is a *kernel* of  $\mathcal{H}$  if

$$S \cap e \cap f \neq \emptyset \quad \forall e, f \in \mathcal{H}.$$

E.g., the kernel of  $\left\{ \binom{S}{r} \right\}$  (for  $|S| < 2r$ ) is itself,  
the kernel of  $K_{1,m}$  (star) is a singleton.

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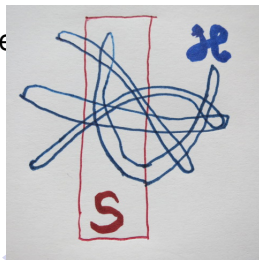
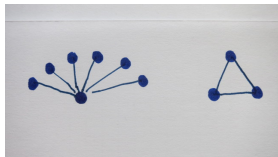
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**Theorem (Całczyńska-Karłowicz 1964)**

$\forall r \exists CK(r) < \infty$  such that:

*If  $\mathcal{H}$  is intersecting,  $|e| \leq r \forall e \in \mathcal{H}$ , then  $\exists S, |S| \leq CK(r)$ .*

$CK(1) = 1$ ,  $CK(2) = 3$ ,  $CK(3) \geq 7$ ,  
 $CK(q+1) \geq q^2 + q + 1$  if  $\exists$  projective plane



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Remove element  $x$  if  $\mathcal{H}|(X \setminus \{x\})$  is still intersecting.

Repeat until we get a critical  $S$ :

$$\forall x \in S \exists A(x), B(x) \in \mathcal{H}_{new} \quad A(x) \cap B(x) = \{x\}.$$

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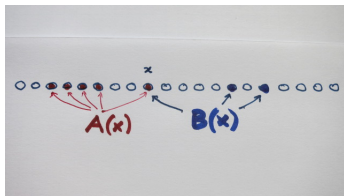


Figure: A type  $x$  permutation of  $X$ .

Call a permutation **type  $x$**  if  $A(x) \leq_{\pi} B(x)$ . Type  $x \neq$  Type  $y$ .  
 $\mathbb{P}(\text{Prob}(\pi \text{ is of type } x)) = \dots, \sum_{x \in S} \mathbb{P}(\text{Prob} \leq 1, \quad CK(r) \leq \frac{r}{2} \binom{2r-1}{r}.$



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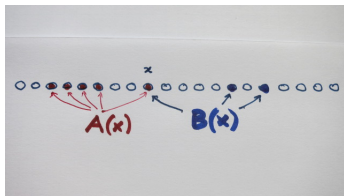


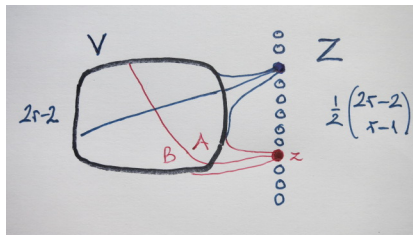
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New results, see: Kang, Ni, Shan 2017, Polcyn, Ruciński 2017, Henning, Yeo 2014, Tuza 1994/1996 (surveys),

# The kernel could be exponential

A **construction** by Erdős, Lovász 1973.

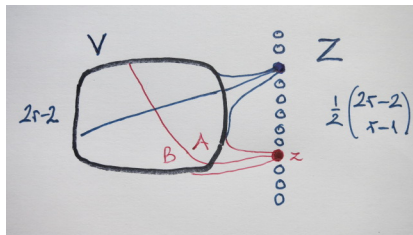


Take  $|V| = 2r - 2$ ,  $|Z| = \frac{1}{2} \binom{2r-2}{r-1}$ ,  $V \cap Z = \emptyset$ .

$\forall$  partition  $\Pi$  of  $V$ ,  $|A| = |B| = r - 1$  assign  $z := z(\Pi) \in Z$ .

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Let  $\{A(z), B(z)\} := \{A \cup \{z\}, B \cup \{z\}\}$ .

Then  $\mathcal{H}$  is intersecting, critical,  $V(\mathcal{H})$  can not be shrunk

$$CK(r) \geq (2r - 2) + \frac{1}{2} \binom{2r - 2}{r - 1} \quad \square$$

New results, generalizations, see: Alon, Füredi 1987, Talbot 2004, Tuza 1994/1996 (surveys), ...

# Sperner families, LYM inequality

DEF:  $\mathcal{H} \subset 2^{[n]}$  is a **Sperner** family if  $A \not\subset B \quad \forall A, B \in \mathcal{H}$ .

Theorem (Sperner)

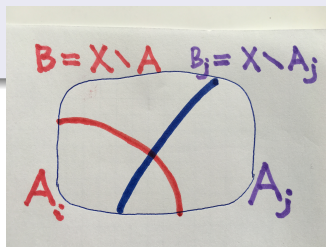
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Theorem (LYM: Lubell, Yamamoto, Meshalkin)

$$\sum_{A \in \mathcal{H}} \frac{1}{\binom{n}{|A|}} \leq 1.$$

Proof:  $(A, X \setminus A)$  is an cross intersecting SPS. □

Equality iff  $|A| = r \quad \forall A \in \mathcal{H}$ .

# Cross intersecting $k$ -tuples

DEF:  $(A_i, B_i, \dots, Z_i)_{1 \leq i \leq m}$  form a family of cross intersecting  $k$ -tuples if these  $k$  sets are pairwise disjoint and

$$\forall i \neq j \quad \exists X \neq Y \in \{A, B, \dots, Z\} \text{ such that } X_i \cap Y_j \neq \emptyset.$$

## Theorem (Tuza)

If  $p_1 + \dots + p_k = 1, \forall p_j > 0$ , then

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Corollary (Take  $p_1 = p_2 = 1/2$ .)

$\{A_i, B_i\}_{1 \leq i \leq m}$  cross intersecting SPS then  $m \leq 2^{\max\{|A_i|+|B_i|\}}$ .

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$\{A_i, B_i\}_{1 \leq i \leq m}$  cross intersecting SPS then  $m \leq 2^{\max\{|A_i|+|B_i|\}}$ .

Even more, it is enough to suppose that for all  $i \neq j$

$$\max\{|A_i \cap B_j|, |A_j \cap B_i|\} > 0.$$

Other results on  $k$ -tuples: Alon 1985.

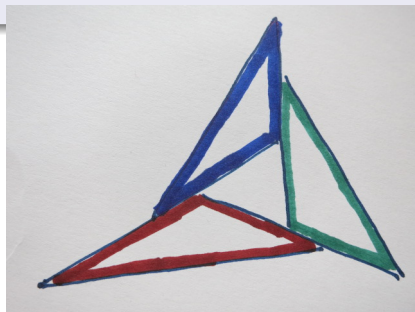


# A geometric application, tangent simplices

DEF:  $S, S' \subset \mathbb{R}^d$  are **tangent** simplices if  
 $\text{int } S \cap \text{int } S' = \emptyset$  and  $\dim(S \cap S') = d - 1$ .

## Theorem (Perles)

Suppose  $S_1, \dots, S_m \subset \mathbb{R}^d$  pairwise tangent.  
Then  $m \leq 2^{d+1}$ .



Four tangent triangles

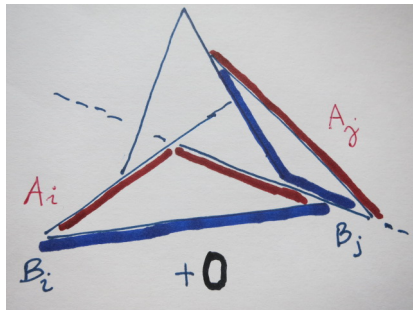
J. Zaks:  $h(2) = 4$ ,  $h(3) = 8$ ,  $h(d) \geq 2^d$ .

# Tangent simplices, a **proof** using SPS

Perles' proof.

Suppose  $S_1, \dots, S_m \subset \mathbb{R}^d$  pairwise tangent.

Let  $H_1, \dots, H_v$  all hyperplanes, tangent to some faces  
( $v \leq m(d+1)$ ).



Take a general point  $O$  and given  $S_i$

$A_i :=$  tangent hyperplanes  $H_\alpha$  such that  $O \in H_\alpha^+$  and  $S_i \subset H_\alpha^+$ .

Then  $\{|A_i \cap B_j|, |A_j \cap B_i|\} = \{0, 1\}$ . □

# Pairs of disjoint subspaces

DEF:  $A, B \subset \mathbb{R}^d$  are **disjoint** linear subspaces if

$$A \cap B = \{\mathbf{0}\}, \quad \text{i.e., } \dim(A + B) = \dim(A) + \dim(B).$$

Non-disjoint:  $\dim(A \cap B) \geq 1$ .

## Theorem (Lovász)

Suppose that  $(A_i, B_i)_{1 \leq i \leq m}$  is a cross intersecting family of disjoint pairs of subspaces of dimensions  $a$  and  $b$ , i.e.,

$$\dim(A_i) \leq a, \dim(B_i) \leq b \text{ and}$$

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Then  $m \leq \binom{a+b}{a}$ .

Implies Bollobás.

Given a cross intersecting SPS  $(\mathcal{A}, \mathcal{B})$ : take  $d = |\cup \mathcal{A} \cup \mathcal{B}|$  pairwise orthogonal vectors  $\{\mathbf{e}_t : t \in [d]\}$  and let  $\hat{A}_i := \text{Span}(\{\mathbf{e}_t : t \in A_i\})$ .

# A consequence of Lovász geometric method

Theorem (ZF **cross  $t$ -intersecting families**)

$(A_1, B_1), \dots, (A_m, B_m)$  set pairs with  $|A_i| \leq a$ ,  $|B_i| \leq b$  and  
 $A_i \cap B_j \leq t$ , and  $A_i \cap B_i > t$  for all  $i \neq j$ .

Then  $m \leq \binom{a+b-2t}{a-t}$ .

Best possible:

$V = X \cup T$ ,  $|X| = a + b - 2t$ ,  $|T| = t$  and

$\mathcal{A} := \{A \in \binom{V}{a}, T \subset A\}$ ,  $\mathcal{B} := \{B \in \binom{V}{b}, T \subset B\}$ ,

# Another consequence, skew cross intersecting SPS's

Theorem (Frankl/ Kalai)

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**Alon and Kalai:** used skew cross intersecting SPS's

to prove the **Upper Bound Theorem of McMullen:**

An exact formula for the max number of  $t$  dim faces of an  $n$ -vertex convex polytope with  $n$  vertices.



# $\tau$ -critical graphs

DEF: (Gallai) The graph  $G$  is  $\tau$ -critical if

$$\tau(G \setminus e) < \tau(G) \quad \forall e \in E(G)$$

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Theorem (Erdős, Gallai)

$G$  is  $\tau$ -critical, then  $|V(G)| \leq 2\tau$ .

Theorem (Erdős, Hajnal, Moon)

$G$  is  $\tau$ -critical, then  $e(G) \leq \binom{\tau+1}{2}$ .

Best possible. (Originally they stated it about  $\overline{G}$ ).

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Best possible. (Originally they stated it about  $\overline{G}$ ).

**Proof** of the EHM thm:

Consider  $A(e) := e$  for edge  $e \in E(G)$ , and

$B(e) :=$  a covering set of the edges of  $E(G) \setminus \{e\}$ .

They form a cross intersecting SPS of set sizes 2 and  $\tau - 1$ . □

# $\tau$ -critical hypergraphs

$$\tau(\mathcal{H} \setminus e) < \tau(\mathcal{H}) \quad \forall e \in E(\mathcal{H})$$

E.g., Erdős-Lovász construction for  $CK(r)$ ,  $K_r^{\tau+r-1}$ .

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Theorem (Petruska, Szemerédi  $r = 3$ , Gyárfás, Lehel, Tuza 1982, finally Tuza)

$\mathcal{H}$  is  $\tau$ -critical, then  $|V(\mathcal{H})| \leq \binom{\tau+r-1}{r-1} + \binom{\tau+r-2}{r-2}$ .

Best upper bound is still unknown.

## PART II. 1-cross intersecting families

DEF: A *1-cross intersecting set pair system* (SPS) of size  $m$ :

$(\{A_i\}_{i=1}^m, \{B_i\}_{i=1}^m)$  with  $A_i \cap B_i = \emptyset$  and  $|A_i \cap B_j| = 1 \quad \forall i \neq j.$



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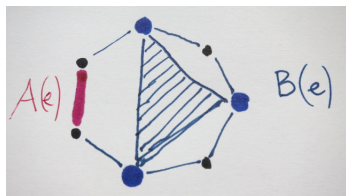
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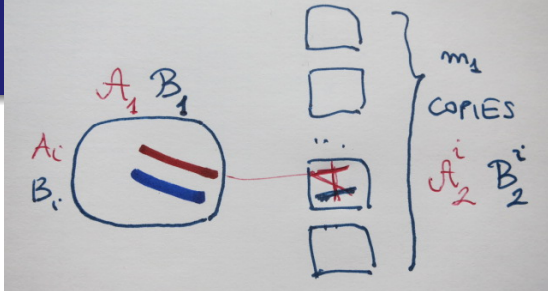
A  $C_5$  and  $\overline{C_5}$  form a  $(2, 2)$ -bounded 1-cross int' SPS.

It is optimal:  $m_2(*, *, 1) = 5$

The edges of a  $C_{2n+1}$  and independent covers:

a  $(2, n)$ -bounded 1-cross intersecting SPS.

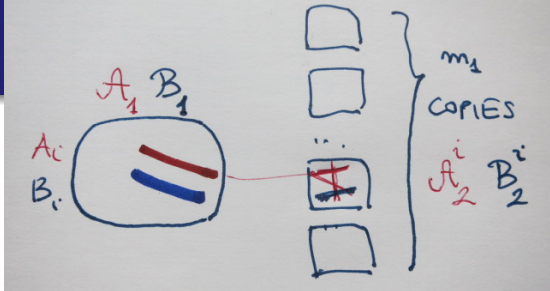
# The product construction



## Proposition

If  $(a_1, b_1)$ -bounded and  $(a_2, b_2)$ -bounded 1-cross intersecting SPS exist with sizes  $m_1$  and  $m_2$ , then  $(a_1 + a_2, b_1 + b_2)$ -bounded 1-cross intersecting SPS also exists of size  $m_1 \cdot m_2$ .

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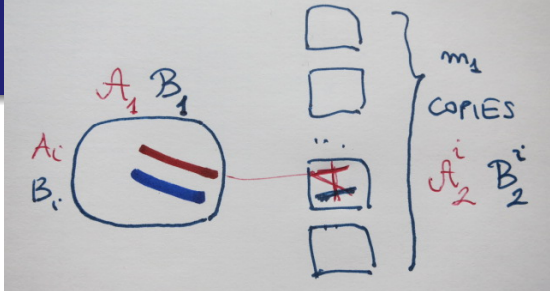
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**Proof** Given  $(\mathcal{A}_1, \mathcal{B}_1)$  and  $(\mathcal{A}_2, \mathcal{B}_2)$  on vertex sets  $V_\alpha$ .

Take  $m_1$  pairwise disjoint ground sets  $V_2^1, \dots, V_2^{m_1}$  with copies  $(\mathcal{A}_2^i, \mathcal{B}_2^i)$ , e.g.,  $\mathcal{A}_2^i = \{A_{i,1}, \dots, A_{i,m_2}\} \dots$

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The pairs  $\mathcal{A}'_{i,j} = \mathcal{A}_i \cup \mathcal{A}_{i,j}$ ,  $\mathcal{B}'_{i,j} = \mathcal{B}_i \cup \mathcal{B}_{i,j}$  form a 1-cross intersecting SPS,  $|\mathcal{A}'_{i,j}| \leq a_1 + a_2$  and  $|\mathcal{B}'_{i,j}| \leq b_1 + b_2$ . □

# 1-cross intersecting SPS can be exponential

## Corollary

*There exists an  $(n, n)$ -bounded 1-cross intersecting SPS of size  $5^{n/2}$  if  $n$  is even and of size  $2 \cdot 5^{(n-1)/2}$  if  $n$  is odd.* □

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The product construction gives a  $(3, 3)$ -bounded 1-cross intersecting SPS of size 10.

We have another example, the pairs  $(\{i, i+1, i+2\}, \{i+3, i+6, i+9\})$  taken (mod 10) has 10 vertices.

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Samuel Spiro (sspiro@ucsd.edu) informed us that his computer program successfully checked that 10 is indeed the largest size,

$$m_1(*, *, 1) = 2,$$

$$m_2(*, *, 1) = 5,$$

$$m_3(*, *, 1) = 10,$$

$$m_4(*, *, 1) \geq 25.$$



Fekete's Lemma on subadditive sequences implies

$$\sqrt{5} \leq \lim_{n \rightarrow \infty} (m_n(*, *, 1))^{1/n} = \exists \leq 4.$$

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A challenging **problem** is to decrease essentially Bollobás' upper bound:

## Conjecture

*There exists a positive  $\varepsilon$  such that  $m_n(*, *, 1) \leq (1 - \varepsilon) \binom{2n}{n}$  for every  $n \geq 2$ .*

# A Fischer's inequality for cross 1-intersecting SPS

Proposition ( $m \leq |\cup \mathcal{A}|$ )

*Let  $(\mathcal{A}, \mathcal{B})$  be 1-cross intersecting,  $V := \cup \mathcal{A}$ . The char. vectors of the edges of  $\mathcal{A}$  are linearly independent in  $\mathbb{R}^V$ .*

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Let  $\mathbf{a}_i$  (resp.  $\mathbf{b}_i$ ) denote the characteristic vector of  $A_i$  (resp.  $B_i$ ),  $\mathbf{a}_i(v) = 1$  **for**  $v \in V$  if and only if  $v \in A_i$ . Otherwise  $\mathbf{a}_i(v) = 0$ .

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Take the dot product with  $\mathbf{b}_j$ . Since  $|A_i \cap B_j| = 1$  for  $i \neq j$  and  $|A_j \cap B_j| = 0$ , we get

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Add up for all  $j$ :  $(m-1) \sum_{i=1}^m \lambda_i = 0$ , consequently  $\sum_{i=1}^m \lambda_i = 0$  and thus  $\lambda_j = 0$  for all  $j$ .



# Lovász' characterization of perfect graphs

A special case of the previous Proposition can be used in the non-trivial part of Gasparyan's 1996 proof of Lovász's theorem:

A graph  $G$  is perfect if and only if

$$|V(H)| \leq \alpha(H)\omega(H) \quad (1)$$

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**Proof.** In a minimal imperfect graph  $G$  there is a 1-cross intersecting SPS of size  $m = \alpha(G)\omega(G) + 1$  defined by independent sets and complete subgraphs. By the previous Proposition,  $|V(G)| \geq \alpha(G)\omega(G) + 1$ , contradicting (1).  $\square$

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## Corollary (Lovász)

*A graph is perfect if and only if its complement is perfect.*

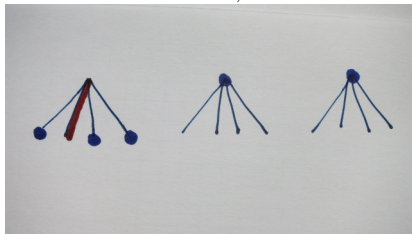
# $(2, n)$ -bounded 1-cross intersecting SPS

Theorem (The case of  $\mathcal{A}$  is a graph.)

Let  $n \geq 4$ , and let  $(\mathcal{A}, \mathcal{B})$  be a  $(2, n)$ -bounded 1-cross intersecting SPS of size  $m$ . Then

$$m \leq \left( \lfloor \frac{n}{2} \rfloor + 1 \right) \left( \lceil \frac{n}{2} \rceil + 1 \right).$$

Best possible. For  $n = 2, 3$  the exact values are  $m = 5, 7$ .



An extremal family:  $\lfloor \frac{n}{2} \rfloor + 1$  copies of stars with  $\lceil \frac{n}{2} \rceil + 1$  edges.

# Linear hypergraphs and cross 1-intersecting SPS

A hypergraph  $\mathcal{H}$  is called *linear* if the intersection of any two different edges has at most one vertex.

E.g., **affine planes**  $AG(2, q)$ . Usually  $V(AG(2, q)) = \mathbb{F}_q^2$ ,  
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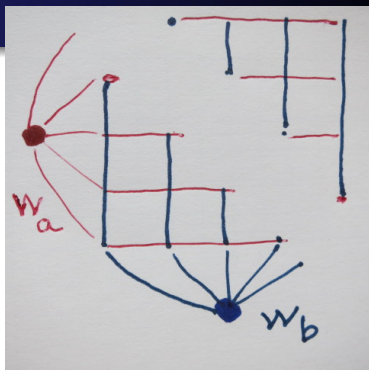
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If one of  $(\mathcal{A}, \mathcal{B})$ , say  $\mathcal{A}$ , in an SPS is linear, then (without any assumption on  $|B_i \cap B_j|, |A_i \cap B_j|$ ).

## Proposition (26)

*The size of an  $(n, n)$ -bounded cross intersecting SPS  $(\mathcal{A}, \mathcal{B})$  with linear  $\mathcal{A}$  is at most  $n^2 + n + 1$ . □*

# Double stars



The vertex set of a *double star of size  $s$*  consist of  $\{v_{i,j} \mid 1 \leq i, j \leq s, i \neq j\}$  and two additional vertices  $w_a$  and  $w_b$ . Define for  $i \in [s]$   $A_i := \{w_a\} \cup \{v_{i,j} \mid 1 \leq j \leq s, j \neq i\}$  and  $B_i := \{w_b\} \cup \{v_{j,i} \mid 1 \leq j \leq s, j \neq i\}$ .

$(\mathcal{A}, \mathcal{B})$  is a **1-cross intersecting** SPS of size  $s$  containing  $s$ -element sets such that **both  $\mathcal{A}$  and  $\mathcal{B}$  are 1-intersecting**.

# Notation and general setting

Let  $a, b > 0$  and  $I_A, I_B, I_{\text{cross}}$  three sets of non-negative integers.

Let  $m(a, b, I_A, I_B, I_{\text{cross}})$  the maximum size  $m$  of a cross intersecting SPS  $(\mathcal{A}, \mathcal{B})$  with the following conditions.

- i)  $A_i \cap B_i = \emptyset$  for every  $1 \leq i \leq m$ ,
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We suppose that  $0 \notin I_{\text{cross}}$ , and  $m \geq 2$ .

If a constraint in iv)–vi) is **vacuous** (i.e.,  $\{0, 1, \dots, |X|\} \subseteq I_X$  or  $\{1, \dots, \min\{a, b\}\} \subseteq I_{\text{cross}}$ ) then we use **the symbol \***

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Bollobás' theorem:

$$m(a, b, *, *, *) = \binom{a+b}{a},$$

# More notations

Our Theorem on p.25 states (for  $n \geq 4$ )

$$m(2, n, *, *, 1) = \left( \left\lfloor \frac{n}{2} \right\rfloor + 1 \right) \left( \left\lceil \frac{n}{2} \right\rceil + 1 \right).$$

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E.g., Proposition p. 26:

$$m_n(01\text{-int}, *, *) \leq n^2 + n + 1.$$

# The origin of the new problems, The *thickness* of a clique partitions

Given a graph  $G$ , a **clique (Biclique) partition** of  $E(G)$   
= parts are complete (complete bipartite) graphs.  
The *thickness* of a partition is the maximum  $s$  such that every  
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**AIM: minimize thickness.**



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# The case of both $\mathcal{A}$ and $\mathcal{B}$ are linear hypergraphs

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Theorem (a bit complicated double counting)

Suppose that  $(\mathcal{A}, \mathcal{B})$  is an  $(n, n)$ -bounded **1-cross intersecting** SPS of size  $m$  such that **both  $\mathcal{A}$  and  $\mathcal{B}$  are linear** hypergraphs. Then  $m \leq \frac{1}{2}n^2 + n + 1$ . I.e.,  $m_n(1\text{-int}, 1\text{-int}, 1) \leq \frac{1}{2}n^2 + n + 1$ .

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Bellow we give three **constructive lower bounds**, large cross intersecting SPS such that  $\mathcal{A}$  is an intersecting linear hypergraph, showing that our results asymptotically the best possible, i.e.,

$$m_n(1\text{-int}, *, 1) \text{ and } \boxed{m_n(1\text{-int}, 1\text{-int}, *) = n^2 - o(n^2)},$$

and  $m_n(1\text{-int}, 1\text{-int}, 1) = \frac{1}{2}n^2 - o(n^2)$ .

(We only give details of the boxed statement.)

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both  $\mathcal{A}$  and  $\mathcal{B}$  are 1-intersecting of sizes  $m_n = n^2 - o(n^2)$ .

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Both  $\mathcal{A}$  and  $\mathcal{B}$  are **almost projective planes!**, and

$A_i \cap B_j = \emptyset \forall i \in [m]$ , and

$A_i \cap B_j \neq \emptyset$  for every  $1 \leq i \neq j \leq m$ .

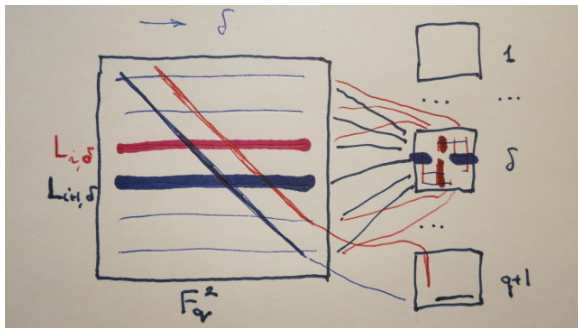
# Second step: The product of an affine plane and the double star

$AG(2, q)$  has  
 $q + 1$  directions  
(parallel classes).  
Each class has  
 $q$  lines.

Let  $\delta$  be a  
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$L_{1,\delta}, \dots, L_{q,\delta}$  the  
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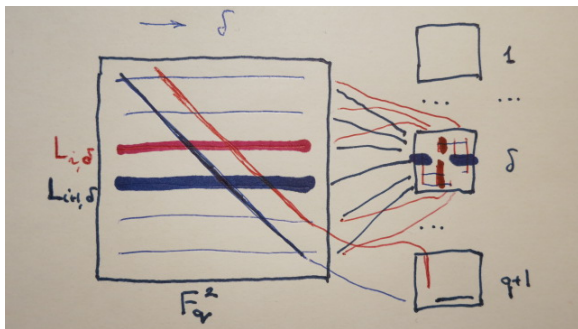
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Let  $A_{i,\delta} := L_{i,\delta} \cup A_i^\delta$  and  $B_{i,\delta} := L_{i+1,\delta} \cup B_i^\delta$ .

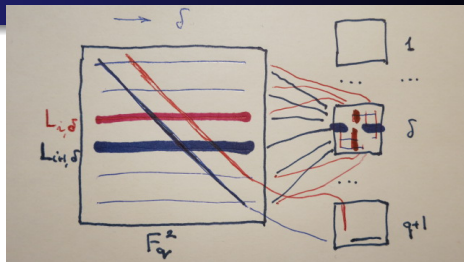
We obtain  $m_{2q} \geq q^2 + q$ .

Call this construction  $\mathcal{H}(2q)$ .



# Cross intersecting almost projective planes

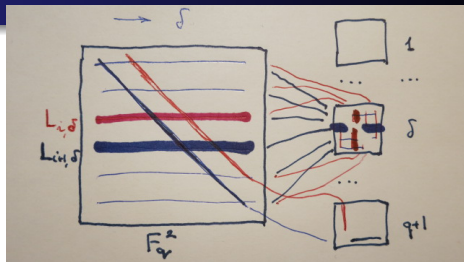
## Third step.



Suppose  $p + 2q \leq n < p + 4q$ , ( $p, q$  primes) where  
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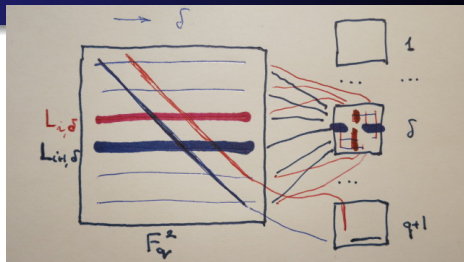


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The size of  $\mathcal{H}(2q)$  is  $q^2 + q \geq p$  thus we need only the first  $p$  set pairs from it. □

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THANKS FOR YOUR ATTENTION!